Modeling and Simulation of Handover Scheme in Integrated EPON-WiMAX Networks

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Abstract

In this paper, we tackle the seamless handover problem in integrated optical wireless networks. Our model applies for the convergence network of EPON and WiMAX and a mobility-aware signaling protocol is proposed. The proposed handover scheme, Integrated Mobility Management Scheme (IMMS), is assisted by enhancing the traditional MPCP signaling protocol, which cooperatively collects mobility information from the front-end wireless network and makes centralized bandwidth allocation decisions in the backhaul optical network. The integrated network architecture and the joint handover scheme are simulated using OPNET modeler. Results show validation of the protocol, i.e., integrated handover scheme gains better network performances.

1. Introduction

During recent years the number of mobile users has grown rapidly. The emerging multimedia services have driven up bandwidth requirements and motivated research of integration and convergence of optical networks and wireless networks.

In the optical regime, Ethernet Passive Optical Network (EPON) has received a worldwide deployment to provide FTTx (x stands for home, building, curb, etc.), because of the bandwidth enhancement and lower maintenance cost offered by optical fibres. In the wireless regime, various wireless access technologies have been developed with an intention to satisfy the continuously growing demand for ubiquitous communication. World Interoperability for Microwave Access (WiMAX) can support diverse communication capabilities for both fixed and mobile users. In the integrated network architecture, the wireline network (EPON) is used to interconnect multiple wireless networks (WiMAX).

Integrated control platform is an important and challenging design issue in the integrated EPON-WiMAX network. In order to ensure an efficient usage of both optical link and radio spectrum, coordinated control could be applied for the overall integrated architecture. In this sense, the roaming among multiple wireless networks, which are the front-end in the integrated network architecture, requires support and assistance from the traditional EPON network. The performance of integrated EPON and WiMAX is studied by designing the traditional and proposed model and conducting discrete event simulations of the models.

There are many studies about handover schemes in wireless networks to provide a fast and seamless roaming. As suggested in previous work, one approach to addressing the mobility issue is to reduce the scanning delay (e.g. channel measurement delay) by either using an efficient scheduling scheme [1] or reducing the number of scanning channels [2]. Another approach focuses on improving handover procedures.

A cross layer handover mechanism is derived in [3]. The combinations of the MAC layer and the network layer messages are utilized to reduce the signaling overhead. The work in [4] aims specifically to improve the handover performances for real-time applications by using a connection ID assignment strategy. Additionally, there are research about joint mobility and network control design, for example, optimal power and handover control design [5] and cell mobility based admission control design [6].

In this paper, a coordinated handover scheme is proposed by enhancing the signaling protocol and the resource allocation scheme in EPON, so that mobility performances in terms of latency and disrupting time are improved in multiple WiMAX networks. The rest of the paper is organized as follows. Section 2 presents the integrated EPON and WiMAX network model. Network design and OPNET simulation are introduced in Section 3. Section 4 describes our proposed mobility aware bandwidth allocation scheme. In section 5 the performance results are analyzed.

2. System Model

We consider the integrated architecture with EPON as the backbone network and a WiMAX network with multiple cells as the front-end network (as shown in Figure 1).

Figure 1: Integrated EPON and WiMAX network architecture
A typical EPON system consists of one Optical Line Terminal (OLT) functionalized as a central office, one passive optical splitter implemented in the remote node, and multiple Optical Network Units (ONUs) residing at subscribers’ locations. The front-end wireless network consists of a set of $k$ cells $\{k = 1, 2, \ldots, K\}$ and a set of $m$ Mobile Users (MSs) $\{m = 1, 2, \ldots, M\}$ in each cell. The service area is partitioned so that each client connects to only one cell at any given time. The handover occurs when a MS travels among cells.

### 2.1 EPON Signaling Control

In the optical domain, the transmission between the OLT and multiple ONUs is based on a poll/grant mechanism. The signaling access protocol in the EPON MAC layer is known as the MultiPoint Control Protocol (MPCP) [7]. MPCP is designed only for the EPON network. There is no mobility concern. In the downstream direction from the OLT to the associated ONUs, data are broadcasted to each ONU in a point-to-multipoint fashion. On the other hand, it is a multipoint-to-point fashion in the upstream direction from the ONUs to the OLT. The OLT allocates upstream bandwidth among the ONUs and each ONU transmits packets in dedicated time slots. The OLT commands and the ONU information are exchanged by two types of MPCP messages: REPORT message is used by an ONU to report bandwidth requirements (normally this value is based on queue occupancies) to the OLT; GATE message is transmitted by the OLT to issue transmission grants to each ONU. The upstream network resource can be distributed among ONUs using either fixed bandwidth allocation algorithms or Dynamic Bandwidth Allocation (DBA) algorithms [8].

### 2.2 WiMAX Mobility

In the wireless domain, WiMAX mobility is based on the IEEE 802.16e standard [9]. In this paper, the hard handover scheme is adopted. The handoff procedure consists of several major stages: network topology acquisition and advertisement, cell scanning, handover decision and initiation, synchronization to the target BS, ranging, and network re-entry.

### 2.3 Proposed Integrated Mobility Management Scheme (IMMS)

The proposed Integrated Mobility Management Scheme (IMMS) utilizes an enhanced MPCP signalling protocol to assist the handover scheme. In order to cooperate with mobility control, the traditional MPCP is extended and new fields are added into the GATE and REPORT signalling messages. We define the control messages deployed in IMMS as $GATE-m$ and $REPORT-m$. The $GATE-m$ and $REPORT-m$ message convey the handover signalling messages and data packets between the OLT and the ONU-BS.

As illustrated in Figure 2, the proposed Integrated Mobility Management Scheme (IMMS) is implemented in the integrated EPON-WiMAX network architecture. When MS moves out of its serving area, a handover to the neighbouring BS is performed to maintain its service connections.

Regarding to the mobility support, the traditional MPCP in EPON has both pros and cons. In the downstream direction, the broadcast ability is of benefit to distribute data to both the serving ONU-BS and the target ONU-BS. When the intention of mobility is reported to the OLT, the data for the MS will be transmitted to both the serving and the target ONU-BS at the same time, which reduces the signalling latency and the packet loss rate during the handover process.

In the upstream direction, ONU-BSs are polled by the OLT in turns and assigned the upstream transmission periods. In the integrated network architecture, multiple ONU-BSs connect to the OLT with no communication among the ONUs. The received mobility control messages and data have to be queued in the OLT and the target ONU-BS until the serving ONU-BS is polled. The traditional poll/request/grant mechanism affects the handover process with additional polling delay. The residue data and the new data after handover are queued until the ONU-BS receives the bandwidth grant from the OLT. Moreover, the polling sequence decides the polling delay, which is the interval between two successive polls to an ONU.

In our proposed IMMS, the signalling scheme is modified to reduce the polling latency, and the additional upstream bandwidth is assigned to provide fast data transmission tunnel for mobile users. We propose new bandwidth allocation schemes in the EPON to provide a fast data transmission tunnel for mobile users. A fixed or dynamic amount of bandwidth, aiming to transmit data for mobile users, is granted to each ONU-BS in addition to the request/grant normal bandwidth.

### 3. OPNET Integrated EPON-WiMAX Simulator

We have developed an integrated EPON-WiMAX model in OPNET (shown in Figure 3). Our system extensively simulates all the essential entities of the network such as OLT, ONU-BS, Network Controller etc. We have considered most of the functionalities at all the protocol layers. In the following sections, we briefly describe the relevant functional blocks of our model implementation. Proposed handover schemes are implemented including mobility and the resource allocation.
3.1 Handover Scheme at ONU-BS

The ONU-BS node is responsible for communications with both the OLT and the MS. Figure 4 shows the basic functionalities including receiving GATE control message (rx_ga state), transmit REPORT message (gen_UL state), communication with MS (rx_ss and tx_DL state), handover operations (cell_up state).

To commence the handover the MS sends mobility-action-intension messages to the serving BS. The MS acquires information about the target network either by receiving network advertisements from the serving BS or scanning of neighbour BSs. In the current IEEE 802.16e specification, the BS can obtain information of neighbouring BSs from the network and periodically broadcasts this information to the MS using a MOB_NBR-ADV message. The MS determines the target BS according to the received network parameters, such as Received Signal Strength Indicator (RSSI) and channel error rate.

A sequence of handover messages are communicated between the serving BS, the target BS and the MS. To establish communication the MS needs to synchronize with the downlink channel of the target BS. The target BS will send the DL/UL MAP to the MS to specify the UL/DL transmission opportunity. The target BS can obtain the MS information from the serving BS through the backbone network. For handover service, the MS performs network re-entry process including reauthorization and re-registration. Afterwards the MS continues with the connection to the target ONU-BSs. Data is received and queued in the target ONU-BS. When the target ONU-BS is polled and assigned with the upstream transmission bandwidth, data is sent.

3.2 Mobility Aware Signaling Mechanism at OLT

Figure 5 shows the OLT process model. The OLToll communicates with ONU-BSs via the GATE and the REPORT messages. Upstream bandwidth is allocated to each ONU-BS to relay data packets. During the handover process, new bandwidth needs to be requested, after the MS moves its connection from the serving ONU-BS 1 to the target ONU-BS k.

4. Enhanced Bandwidth Allocation Scheme

The OLT communicates with ONU-BSs via the GATE and the REPORT messages. Upstream bandwidth is allocated to each ONU-BS to relay data packets. During the handover process, new bandwidth needs to be requested, after the MS moves its connection from the serving ONU-BS 1 to the target ONU-BS k.

4.1 Bandwidth Allocation Problem Statement

The bandwidth allocation for the upstream transmission is illustrated in Figure 6. At \( t_0 \), MS sends the mobility intension to ONU-BS 1, which relays the handover request to the OLT via Report1 message at \( t_1 \). When the OLT polls the ONU-BS k, the mobility request message is embedded and transmitted with the Gatek message at \( t_2 \). The response from the target ONU-BS k is forwarded back to ONU-BS 1, when ONU-BS 1 is polled again after a polling cycle \( (T_{polling} = t_3 - t_0) \). The polling cycle is the interval between two successive polls to an ONU-BS. During a polling cycle, all connected ONU-BSs are polled and finished their transmission with the OLT. The length of a polling cycle is determined by the total number of ONU-BSs and their granted transmission periods. If the handover intension is accepted, the reside data (\( d_{res} \)) in the serving ONU-BS 1 needs to be transferred to the target ONU-BS k. However, the data cannot be sent to the OLT immediately. The serving ONU-BS 1 should
request for upstream bandwidth (at $t_1$). The granted bandwidth is determined by the OLT and is assigned by the GATE message in the next polling cycle (at $t_2$). Similar to the residual data, after the MS re-registers to the target ONU-BS (at $t_3$), the new arrival traffic ($d_{\text{new}}$) has to wait in the queue (at $t_4$) until the ONU-BS receives the assigned time slot (at $t_5$).

Using the traditional MPCP, it is noted that a complete data transmission experiences additional handover delay, which is caused by the polling procedures in EPON. The polling cycle ($T_{\text{polling}}$) is calculated as the sum of all granted bandwidth ($BW_{\text{grant}}$) to $K$ ONU-BSs (illustrated in Eq (1)). $R_o$ is the transmission rate of the optical uplink and $T_g$ is the guard time between two successive upstream data.

$$T_{\text{polling}} = \sum_{k=1}^{K} (\frac{BW_k}{R_o} + T_g) \tag{1}$$

The values of delay to transmit the residue data ($D_{\text{res}}$) and the new arrival data ($D_{\text{new}}$) are calculated in Eq (2) and Eq (3).

$$D_{\text{res}} = t_4 - t_0 = 2 \times T_{\text{polling}} + (t_1 - t_0) \tag{2}$$

$$D_{\text{new}} = t_5 - t_3 = T_{\text{polling}} + (t_6 - t_3) \tag{3}$$

There is a polling delay occurred after the MS indicates the mobility intention to the OLT and when the MS connects to the new ONU-BS at $t_5$, which increases the handover latency. Due to the lack of interoperability and support of mobility in EPON, this additional delay is inevitable.

### 4.2 Fixed Mobile Bandwidth Grant (FMBG) scheme

In this first strategy, the OLT schedules a fixed amount of bandwidth for each ONU-BS, aiming to transmit handover data. This is a simple and direct solution to eliminate the polling delay. The total bandwidth ($BW_{\text{total}}$) allocated by the OLT to each ONU-BS consists of the requested bandwidth ($BW_{\text{grant}}$) and the fixed amount of bandwidth for handover usage ($BW_{\text{ho}}$). In the FMBG scheme, buffered data and new data can be transmitted during the additional assigned time slots. After receiving response from the target ONU-BS $k$ at $t_3$, the serving ONU-BS 1 uploads the residue data to the OLT immediately using the additional allocated bandwidth.

Similarly, MS continues the communication with the OLT when it shifts to the new cell (at $t_5$). Therefore, the delay of requesting and granting bandwidth during handoff is avoided. However, the value of the dedicated bandwidth to the handover is difficult to be determined. Assigning a small value may not be able to transmit the buffered data and assigning a large value may increase the interval of polling cycle. The calculation of a polling cycle and the delay for the residue data and new arrival data are shown in Eq (4)-(6).

$$T_{\text{polling}} = \sum_{k=1}^{K} (\frac{BW_{\text{grant}}}{R_o} + BW_{\text{ho}}/R_o) \quad \tag{4}$$

$$D_{\text{res}} = T_{\text{polling}} + (t_1 - t_0) \quad \tag{5}$$

$$D_{\text{new}} = t_0 - t_5 \quad \tag{6}$$

Compared to the traditional scheme, the FMBG scheme supports the handover by simply assigning more bandwidth. Especially the new arrival data can be transferred shortly without causing a considerable pause during conversation. However, the special bandwidth for handover introduces increased polling cycle and polling delay for each ONU-BS.

### 4.3 Dynamic Mobile Bandwidth Grant (DMBG) scheme

The second proposed strategy, the additional bandwidth for handover process is allocated to the ONU-BS, which sends mobility requests. Compared with the traditional scheme, the serving ONU-BS sends the mobility intention request to the OLT at $t_1$. Note that at the same time ($t_1$), the size of buffered data ($d_{\text{buf}}$) will be indicated via the REPORT message. The OLT allocates the request bandwidth when the OLT poll the serving ONU-BS next time. Differed from the traditional scheme is that the residue data will be transmitted at $t_5$.

After the response from the target ONU-BS is returned (at $t_3$), the OLT assigns an approximate amount of the upstream bandwidth to the target ONU-BS, which is aimed to transmit data for the new connected mobile users. The value of bandwidth is an estimation based on the previous transmission of mobile users and provided by the serving ONU-BS. When the re-registration with the new cell is completed at $t_5$, the target ONU-BS will receive additional bandwidth and new arrival data can be uploaded at $t_6$. The handover delay is reduced, because the polling period for new connected MS is eliminated.

Compared to the traditional scheme, the dedicated bandwidth is
assigned to the serving and the target ONU-BS in order to reduce the handover latency. Compared to the FMBG scheme, the additional bandwidth is allocated only to the cell with mobility intension, which eliminates the delay for all ONU-BS. However, information (such as buffered data size) is required from the serving ONU-BS, which results in complex control and more overhead. The value of polling cycle changes when there is additional bandwidth granted for the handover process.

5. Performance Evaluation

We simulated the integrated mobility scheme in the integrated EPON and WiMAX network architecture using OPNET [10]. Network topology is shown in Figure 1. The simulation parameters are shown in Table 1 and Table 2. Traffic arrival process follows a Poisson distribution and the mean arrival rate varies from 0.01 to 0.1 Mbps per MS. For simplicity there is one ONU-BS per wireless cell and MSs are distributed randomly over the cell region.

![Figure 6: Traditional MPCP with handover data transmission Performance results: handover latency vs. traffic load;](image)

Figure 6: Traditional MPCP with handover data transmission Performance results: handover latency vs. traffic load.

![Figure 7: Traditional MPCP with handover data transmission Performance results: handover latency under different number of ONU-BSs.](image)

Figure 6: Traditional MPCP with handover data transmission Performance results: handover latency under different number of ONU-BSs.

6. Conclusion

In the integrated EPON and WiMAX network, a seamless handover and Quality of Service (QoS) provision is an important issue. In this work, we analyze the problem in supporting the handover process and propose two enhanced scheme. The scheme can avoid the additional polling delay caused in the optical domain. The integrated network architecture and functionalities are modelled in OPNET. Extensive simulations are carried out based on different traffic profiles. By performance evaluation, we show that the proposed scheme can reduce the handover latency.

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References


